



Introduction



Non-Problem



Usage



Implementation



Conclusion

Object-Oriented Approaches to Programming

～ Case Study: Binary Methods ～

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Types, Classes, Inheritance

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- ▶ **Binary Operation:** 2 arguments of the same type
Arithmetics, ordering relations ($=$, $+$, $>$ etc.)
- ▶ **OO Programming:** 2 objects of the same class
Benefit from polymorphism
- ▶ Hence the expression “binary method”
- ▶ **[Bruce et al., 1995]:**
 - ▶ Problematic concept in traditional OO approach
 - ▶ Type / class relation in the context of inheritance

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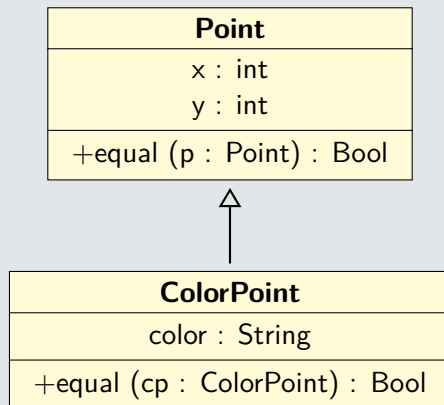
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The Point Hierarchy



C++: Attempt #1

```
class Point
{
    int x, y;

    bool equal (Point& p)
    { return x == p.x && y == p.y; }
};

class ColorPoint : public Point
{
    std::string color;

    bool equal (ColorPoint& cp)
    { return color == cp.color && Point::equal (cp); }
};
```

C++: Attempt #1

Fail

```
int main (int argc, char *argv[])
{
    Point& p1 = * new ColorPoint (1, 2, "red");
    Point& p2 = * new ColorPoint (1, 2, "green");

    std::cout << p1.equal (p2) << std::endl;
    // => True. #### Wrong !
}
```

- ▶ `ColorPoint::equal` masks `Point::equal` (statically)
- ▶ Only `Point::equal` is seen from the base class
- ▶ We need the exact version

C++: Attempt #2

```
class Point
{
    int x, y;

    virtual bool equal (Point& p)
    { return x == p.x && y == p.y; }
};

class ColorPoint : public Point
{
    std::string color;

    virtual bool equal (ColorPoint& cp)
    { return color == cp.color && Point::equal (cp); }
};
```

C++: Attempt #2

A ColorPoint

A Point

Fail

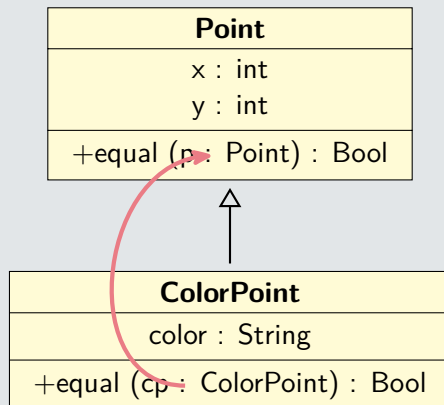
```
bool foo (Point& p1, Point& p2)
{
    return p1.equal (p2);
}
```

ColorPoint::equal
expects a ColorPoint...

...but only gets
a Point !

Context

The Point Hierarchy



Reminder: Covariance / Contravariance

▶ Theoretic Rule:

- ▶ *contravariance* on arguments
- ▶ *covariance* on return values

```
virtual bool equal (Point& p);  
virtual bool equal (ColorPoint& cp);
```

▶ In Practice:

- ▶ *invariance* on arguments
- ▶ Ambiguities due to overloading

▶ Remark:

- ▶ Eiffel allows covariance on arguments
- ▶ Potential errors at run-time

▶ Analyse: [Castagna, 1995]

- ▶ Lack of expressivity
sub-typing (by sub-classing) $\neq$ specialization
- ▶ Defect of the traditional OO model
Lack of multi-methods

CLOS: Attempt #1

```
(defgeneric point= (a b))

(defclass point ()
  ((x :initarg :x :reader point-x)
   (y :initarg :y :reader point-y)))

(defmethod point= ((a point) (b point))
  (and (= (point-x a) (point-x b))
        (= (point-y a) (point-y b))))

(defclass color-point (point)
  ((color :initarg :color :reader point-color)))

(defmethod point= ((a color-point) (b color-point))
  (and (string= (point-color a) (point-color b))
        (call-next-method)))
```

CLOS: Attempt #1

Success

```
(let ((p1 (make-point :x 1 :y 2))
      (p2 (make-point :x 1 :y 2))
      (cp1 (make-color-point :x 1 :y 2 :color "red"))
      (cp2 (make-color-point :x 1 :y 2 :color "green"))))
  (values (point= p1 p2)
          (point= cp1 cp2)))
;; => (T NIL)
```

- ▶ Selection of the appropriate method based on the 2 arguments *Multiple dispatch*
- ▶ Call syntax more pleasing
p1.equal(p2) or p2.equal(p1)?
- ▶ **“Binary Function”**

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Reminder: Method Combinations

▶ Avoid code duplication:

- ▶ C++: `Point::equal()`
- ▶ CLOS: `(call-next-method)`

▶ Applicable methods:

- ▶ All methods compatible with the arguments' classes
- ▶ Sorted by decreasing specificity order
- ▶ `call-next-method` executes the next method in that ordering

▶ Method combinations:

- ▶ Call several applicable methods
Not only the most specific
- ▶ Predefined combinations: `and`, `or`, `progn` etc.
- ▶ Programmable

The and Combination

```
(defgeneric point= (a b)
  (:method-combination and))

(defmethod point= and ((a point) (b point))
  (and (= (point-x a) (point-x b))
        (= (point-y a) (point-y b))))

(defmethod point= and ((a color-point) (b color-point))
  (and (call-next-method)
    (string= (point-color a) (point-color b)))))
```

► **Note:** the CLOS dispatch is programmable

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Non-Conforming Usage

Illegitimate Success

```
(let ((p (make-point :x 1 :y 2))  
      (cp (make-color-point :x 1 :y 2 :color "red")))  
  (point= p cp))  
;; => T #### Wrong !
```

- ▶ (point= <point> <point>) is applicable
A color-point is a point
- ▶ We want to avoid that situation

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Introspection in CLOS

```
(assert (eq (class-of a) (class-of b)))
```

► When to Check ?

- Every method: duplication
- Last method: not efficient / not always possible
- :before: unavailable outside the standard combination
- New combination: not the place
- :around: yes but...

► Where to Check ? (better question)

- Nowhere near the code for point=
- Related to the concept of binary function, not to point=
- Express the concept of binary function itself
 - *Binary functions are special kinds of generic functions*

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A Binary Function Meta-Class

- ▶ Generic functions: `standard-generic-function`
- ▶ Function-objects: `funcallable-standard-class`

```
(defclass binary-function (standard-generic-function)
  ()
  (:metaclass funcallable-standard-class))

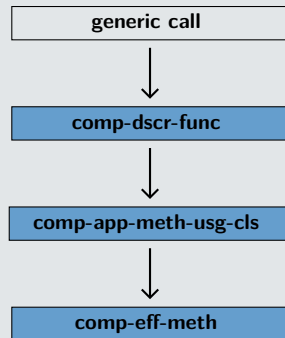
(defmacro defbinary (function-name lambda-list &rest options)
  `(defgeneric ,function-name ,lambda-list
    (:generic-function-class binary-function)
    ,@options))
```

- ▶ **Remark:** the concept mandates a specific `lambda-list`...

Generic Calls

- ▶ `compute-discriminating-function`:
the generic function's actual function
- ▶ `compute-applicable-methods-using-classes`:
figure out which methods are applicable based
on the arguments' classes
- ▶ `compute-effective-method`:
create the (combined) global method for these
arguments
- ▶ **Note:** *specializable* generic functions

Protocol



Application to Binary Functions

```
(defmethod compute-applicable-methods-using-classes
  :before ((binary-function binary-function) classes)
  (assert (= (length classes) 2))
  (assert (apply #'eq classes)))
```

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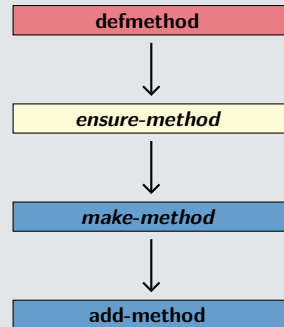
```
;; #### WRONG SPEC!!  
(defmethod point= ((p1 point-3d) (p2 point-2d))  
  #/ ... /#)
```

- ▶ Previously: protect against illegal calls
- ▶ New question: protect against illegal implementations ?

Methods Definitions

- ▶ `ensure-method`:
creation and addition of a new method
- ▶ `make-method`:
creation of the method meta-object
- ▶ `add-method`:
addition and update to the generic function
- ▶ **Note:** *specializable* generic functions

Protocole



Application to Binary Functions

```
(defmethod add-method :before ((bf binary-function) method)
  (let ((method-specializers (method-specializers method)))
    (assert (= (length method-specializers) 2))
    (assert (apply #'eq method-specializers))))
```

- ▶ method: meta-object (Cf. standard-method)
- ▶ method-specializers: accessor

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Missing Implementation

- ▶ Every sub-class of point must specialize point=
- ▶ Late checking: at generic call time
preserve high dynamicity
- ▶ “Binary-Completeness”: there exists a primary method applicable for every class in the hierarchy

```
(defmethod compute-applicable-methods-using-classes
  :around ((bf binary-function) classes)
  (multiple-value-bind (methods ok) (call-next-method)
    ;; 1. Check for binary completeness on METHODS
    ;; 2. If everything goes well, simply return
    ;;    what we got from CALL-NEXT-METHOD
    (values methods ok)))
```

Implementation

```
(loop :for class :in (class-precedence-list (car classes))
      :until (eq class (find-class 'standard-object))
      :do (assert
            (find-if
             (lambda (method)
               (let ((qualifiers (method-qualifiers method)))
                 (and (equal (method-specializers method)
                             (list class class))
                      (not (member :around qualifiers))
                      (not (member :before qualifiers))
                      (not (member :after qualifiers))))))
            methods)))
```

- ▶ class-precedence-list: accessor
- ▶ method-qualifiers: accessor

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- ▶ **Binary methods:** a very simple concept
 - ▶ Difficult to implement in traditional OO systems
 - ▶ Trivial with multi-methods
- ▶ **Thanks to CLOS:**
 - ▶ Usage validation
 - ▶ Implementation validation
- ▶ **Thanks to the MOP:**
 - ▶ Reified Concept
 - ▶ Extension of the core OO system



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Bibliography

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